

# Adaptive Multipath Routing for Dynamic Traffic Engineering

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## Outline

- Introduction to traffic engineering
- Adaptive Multi-Path (AMP) algorithm
- Performance evaluation and results
- Summary and outlook

## What is "Traffic Engineering" (TE)?

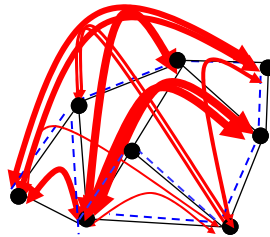


- Traffic engineering is defined as performance optimization of operational networks (IETF)
  - Consider the traffic at the macroscopic level
  - Consider the network as a set of *limited* resources
    - Transmission bandwidth, switching throughput
- Traffic engineering tries to optimally match **traffic demands** with the available **network** resources by acting on **routing**

**Traffic Demands**

**Network**

**Routing**



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## What can TE do, what not



- Objective: To balance the load in the network
  - Prevent mid-term congestion in the core
  - Increase the "effective capacity" of the network
  - Adapt to macroscopic changes in the traffic distribution (time scale: several minutes and above)
- What TE does not do:
  - Counteract congestion in the access links
  - Introduce service differentiation and/or guarantees (→ job for QoS mechanisms)
  - React to faults in the short-term (→ job for Resilience mechanisms)

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- By combating congestion TE indirectly improves the QoS
  - but TE is not a QoS mechanism
  - TE and QoS act at different levels and time-scales
  - Possible interaction: differentiated TE per QoS class
- TE hold strict relationships with resilience techniques
  - Fault isolation implies diverting traffic to another path

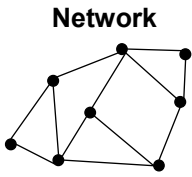


- Traffic engineering methods for IP networks:
  - Link weight optimization in native IP networks
  - Optimization of Multi-Protocol Label Switched (MPLS) networks
  - Algorithmic approaches (dynamic routing in the ARPANet, OMP)

# Example of Connection-Less TE: Link Weight Optimization

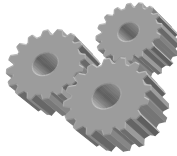


Traffic Demands

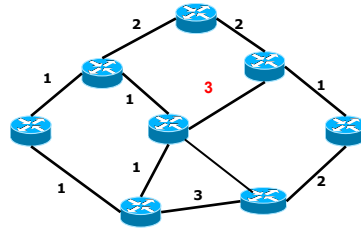


Network

Optimization..



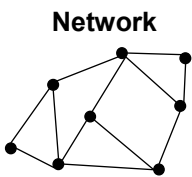
Set of Link Weights



# Example of Connection-Oriented TE: Explicit-Routing Optimization

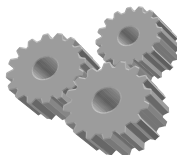


Traffic Demands

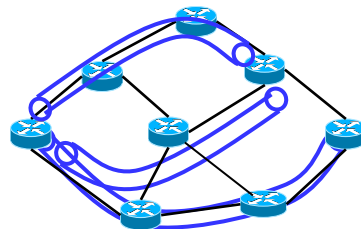


Network

Optimization..

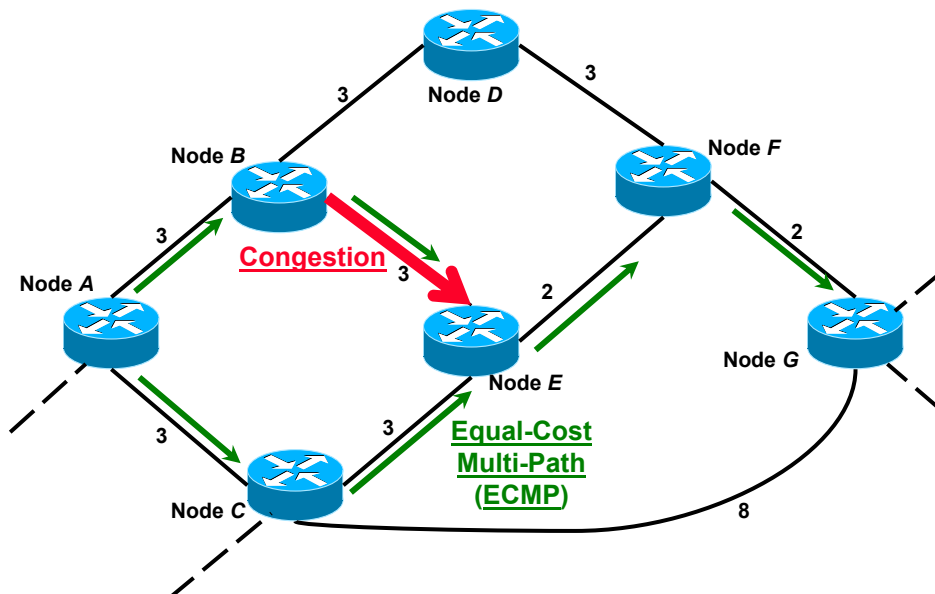


Set of Explicit Routes  
for Virtual Pipes

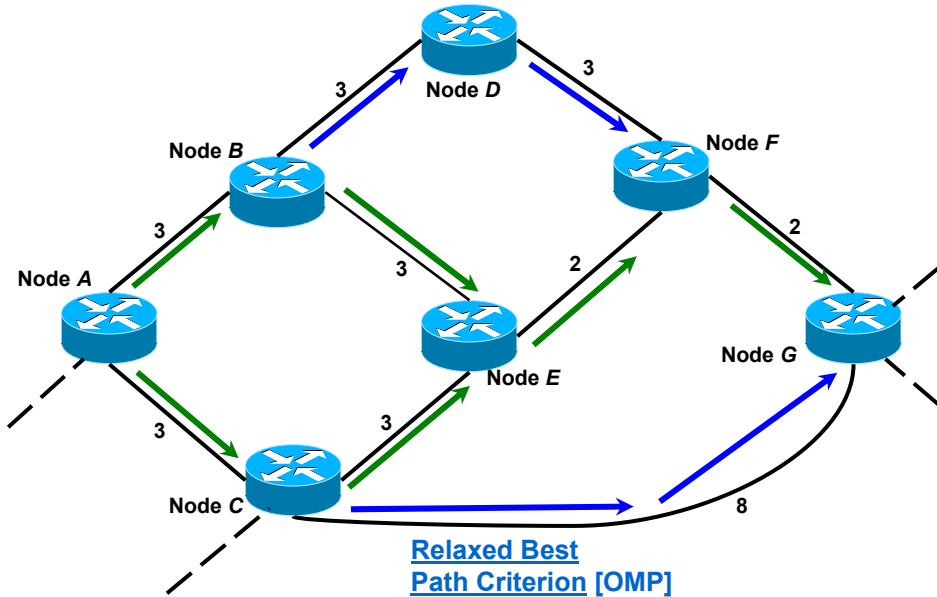


- Existing traffic engineering methods have important disadvantages:
  - MPLS and *link weight optimization* require additional network management
  - Unpredictable signaling overhead with Optimized Multi-Path (OMP)
  
- Our objective:
  - Autonomous and continuous load distribution in the network
  - Low overhead in terms of memory and bandwidth consumption
  
- Proposal: Adaptive Multi-Path Algorithm (AMP)

## Adaptive Multi-Path (AMP)



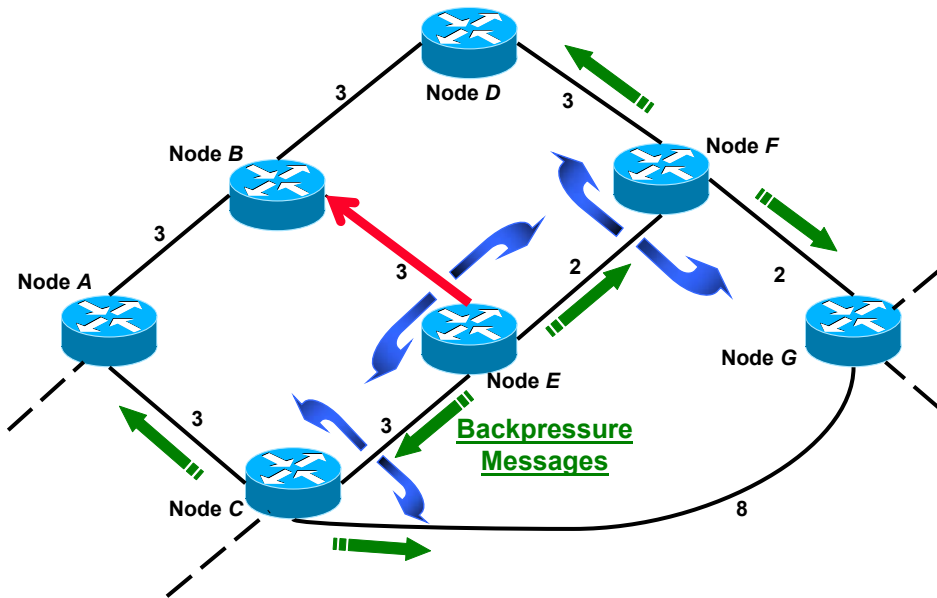
# AMP – Choice of Multiple Paths



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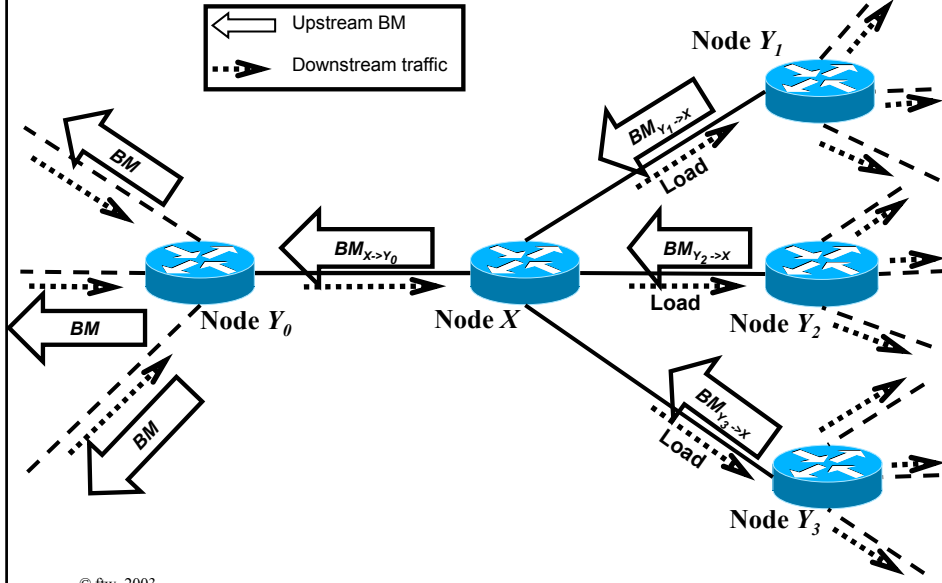
# AMP – Basic Operation



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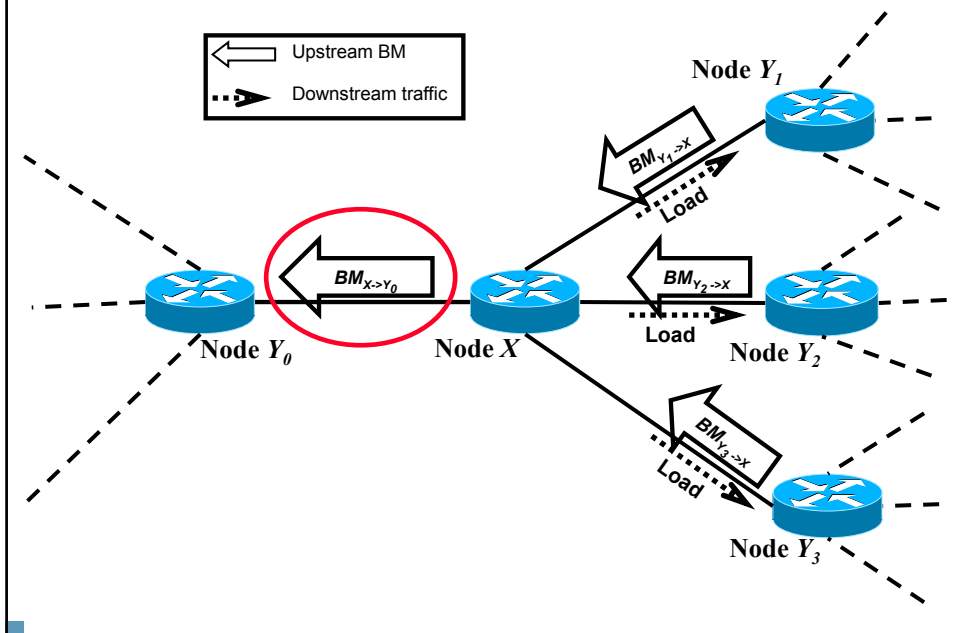
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# AMP – Signaling



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# AMP – Signaling



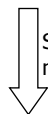
$$BM_{X \rightarrow Y_0} = f(\text{Load}_{XY_1}, \dots, \text{Load}_{XY_n}, BM_{Y_1 \rightarrow X}, \dots, BM_{Y_n \rightarrow X})$$

**Quasi-recursive structure of  
backpressure messages**

⇒

**GLOBAL PROPAGATION OF LOAD  
INFORMATION THROUGH LOCAL  
EXCHANGE OF SIGNALING MESSAGES**

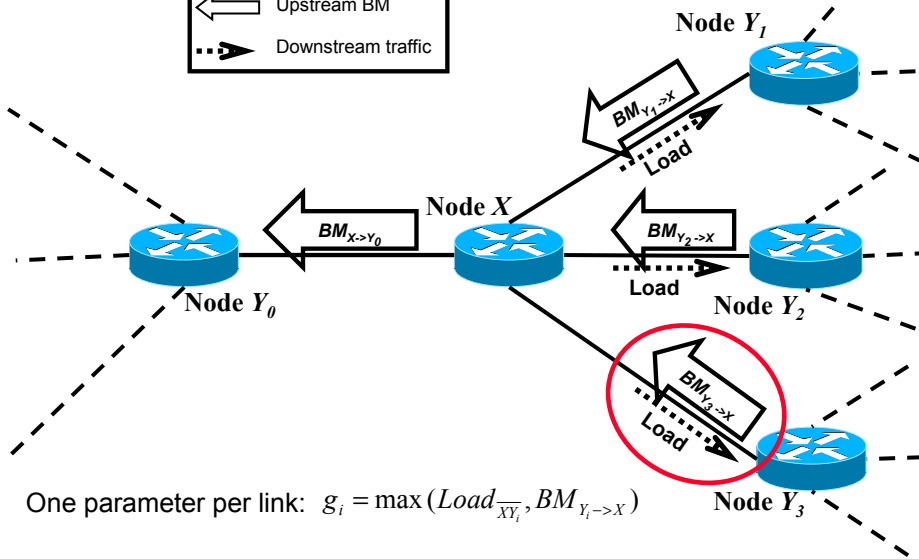
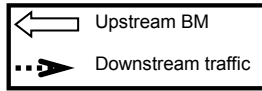
$$BM_{X \rightarrow Y_0} = f(\text{Load}_{XY_1}, \dots, \text{Load}_{XY_n}, BM_{Y_1 \rightarrow X}, \dots, BM_{Y_n \rightarrow X})$$



Summarization of the  
number of parameters



# AMP – Signaling



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# AMP – Signaling



$$BM_{X \rightarrow Y_0} = f(\text{Load}_{XY_1}, \dots, \text{Load}_{XY_n}, BM_{Y_1 \rightarrow X}, \dots, BM_{Y_n \rightarrow X})$$

Reduction of the number of parameters

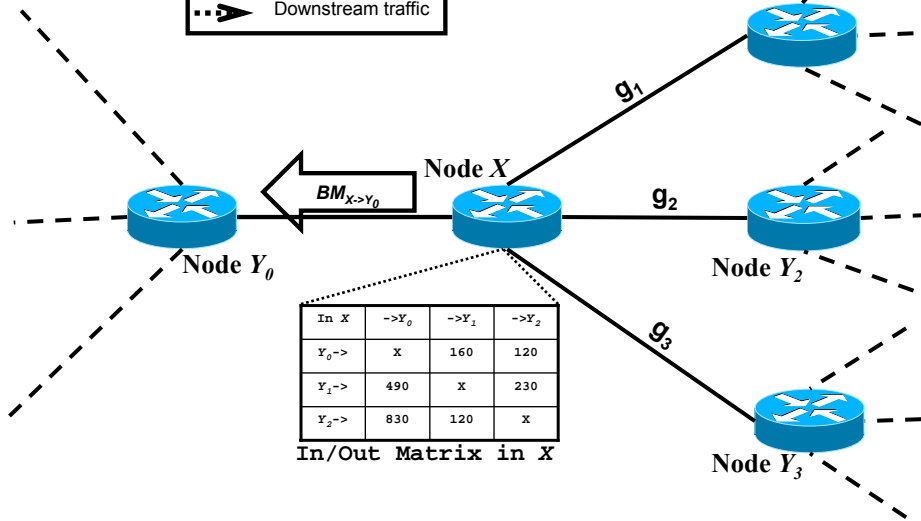
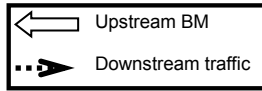
$$g_i = \max(\text{Load}_{XY_i}, BM_{Y_i \rightarrow X})$$

$$BM_{X \rightarrow Y_0} = f(g_1, g_2, \dots, g_n)$$

?

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# AMP – Signaling



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# AMP – Signaling



$$BM_{X \rightarrow Y_0} = f(\text{Load}_{XY_1}, \dots, \text{Load}_{XY_n}, BM_{Y_1 \rightarrow X}, \dots, BM_{Y_n \rightarrow X})$$

Reduction of the number of parameters

$$g_i = \max(\text{Load}_{XY_i}, BM_{Y_i \rightarrow X})$$

↓

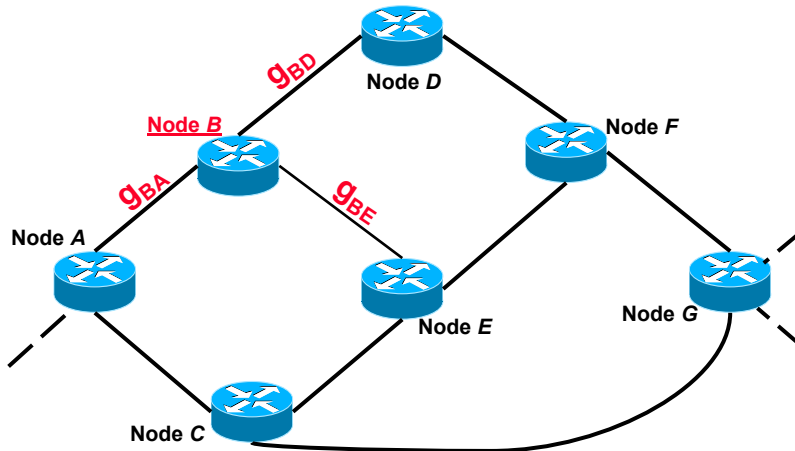
$$BM_{X \rightarrow Y_0} = f(g_1, g_2, \dots, g_n)$$

$$= \sum_{Y_i \in \Omega_X \setminus Y_0} \frac{\beta_{Y_0 X Y_i}}{\beta_{X Y_i}} \cdot g_i$$

weights for congestion contributions

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## AMP – Load Balancing



- The goal of the load balancing mechanism in every node is to equalize the values of  $g$  on all output links

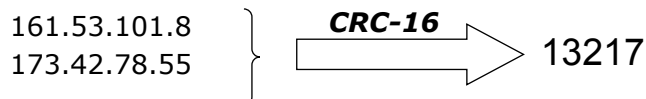
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## AMP – Load Balancing



- In order to avoid packet disordering:
  - => the unit for load balancing is a microflow aggregate
  - => packets are assigned to an aggregate by applying a *CRC-16* hash-function on their source and destination IP addresses
- The *CRC-16* solution space  $[0, 65535]$  is divided among the viable next hops



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## AMP – Load Balancing



- Example routing table in Node *B* – the hash-space boundaries are defined for every reachable destination

Destinations (in Node <i>B</i> )	Next hop: Node <i>A</i>	Next hop: Node <i>D</i>	Next hop: Node <i>E</i>
Node <i>A</i>	[0 – 65535] (ALL PACKETS)		
Node <i>C</i>	[0 – 23723]		[23724 – 65535]
Node <i>D</i>		[0 – 65535] (ALL PACKETS)	
Node <i>E</i>			[0 – 65535] (ALL PACKETS)
Node <i>F</i>		[0 – 34447]	[34448 – 65535]
Node <i>G</i>		[0 – 52142]	[52143 – 65535]

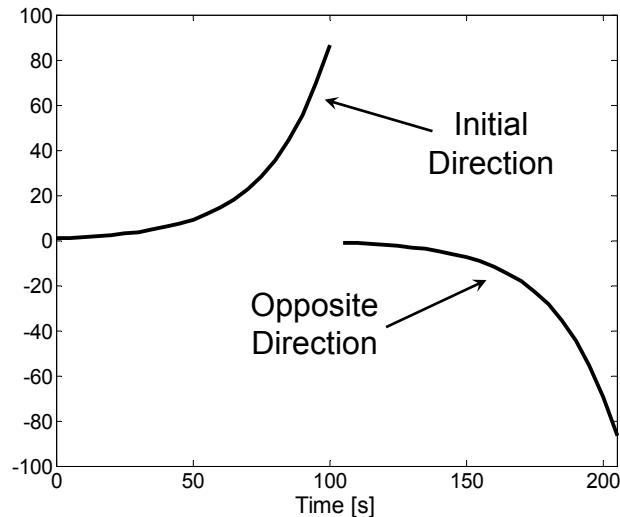
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## AMP – Load Balancing

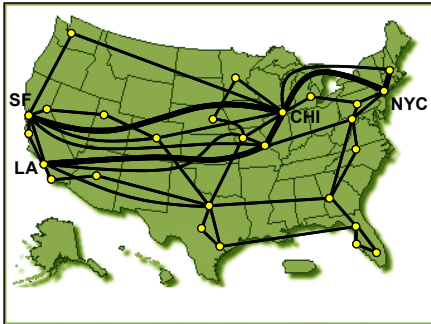


- Conservative load balancing mechanism – the size of load adjustment steps is changed dynamically

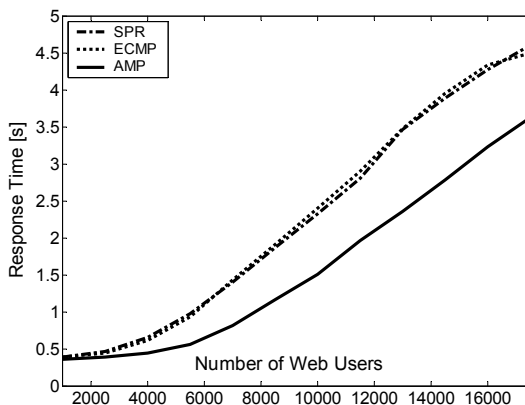


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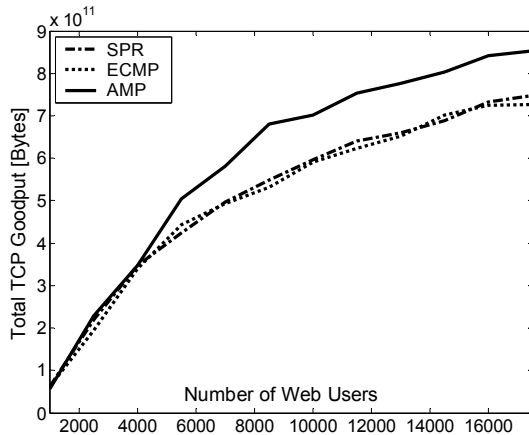
- Implementation of AMP in Network Simulator (ns-2)
- Simulated topology:
  - AT&T-US Network of 27 nodes and 47 links
  - Link capacities of 2.4 and 9.6 Gbit/s (scaled down to 15 and 60 Mbit/s in our simulations)
- Simulated traffic:
  - Web traffic according SURGE model
  - Traffic distribution according to the gravity model
  - Linear scaling of the number of Web users



- Web page response time most important metric from the user's perspective
- Significant reductions in Web page response times throughout investigated scenarios (up to 43%)

- SPR – Shortest Path Routing
- ECMP – Equal-Cost Multi-Path Routing

## AMP Performance Evaluation – Total TCP Goodput

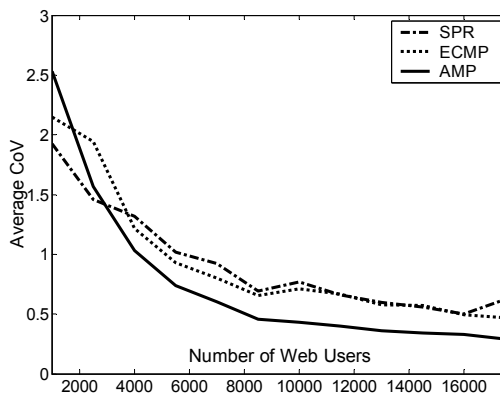


- Improved efficiency of resource utilization
- Total TCP goodput consistently higher with AMP compared to SPR and ECMP in our simulations (improvements of up to 28%)

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## AMP Performance Evaluation – Average CoVs of Link Load



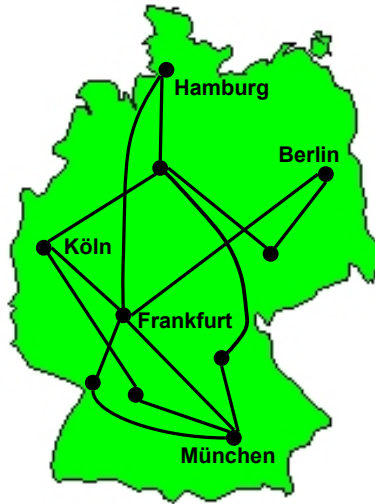
- Similar average Coefficient of Variation (CoVs) of all link loads for the three routing strategies

⇒ stability of AMP load balancing

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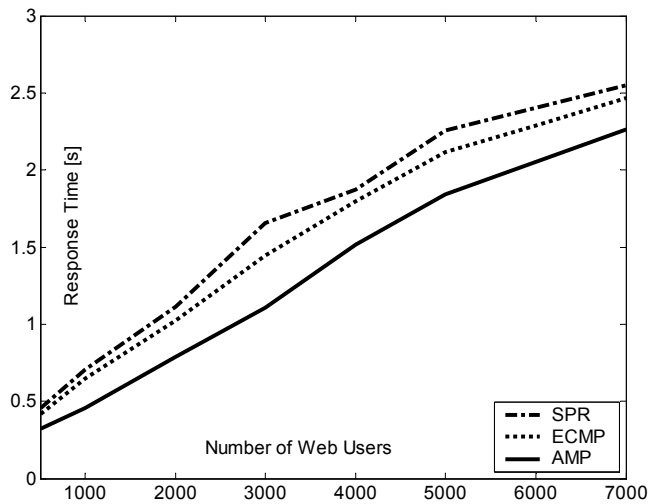
# AMP Performance Evaluation



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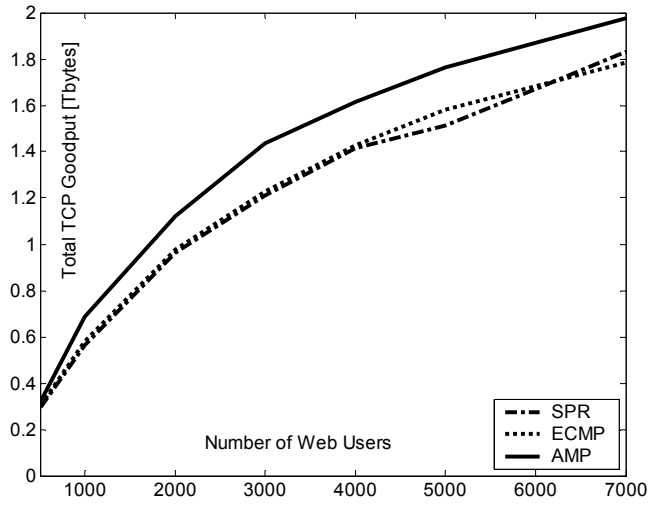
# AMP Performance Evaluation – Average Web Page Response Time



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## AMP Performance Evaluation – Total TCP Goodput



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## Publications



- I. Gojmerac, T. Ziegler, P. Reichl: *Adaptive Multipath Routing Based on Local Distribution of Link Load Information*. Proc. QoFIS'03, Stockholm, October 2003.
- I. Gojmerac, T. Ziegler, F. Ricciato, P. Reichl: *Adaptive Multipath Routing for Dynamic Traffic Engineering*. Proc. GLOBECOM'03, San Francisco, November 2003.

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- AMP Summary:
  - Load balancing within the framework of routing
  - No management overhead, minimal signaling overhead
  - Implementation in Network Simulator (ns-2)
  - Significant performance improvements
- Future research:
  - AMP and network resilience

**Thank you for  
your attention!**